

Analysis of Algorithms, I

CSOR W4231.002

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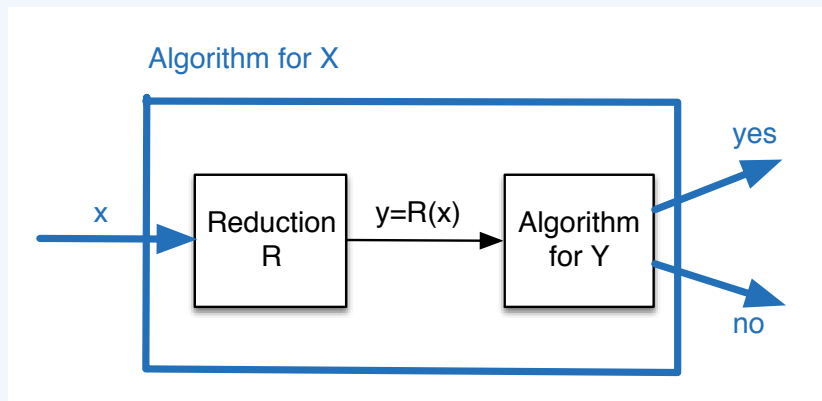
Thursday, April 7, 2016

- 1 Review of last lecture
 - The class \mathcal{NP}
 - The class of \mathcal{NP} -complete problems
- 2 *Satisfiability*: a fundamental \mathcal{NP} -complete problem
- 3 The art of proving \mathcal{NP} -completeness
 - $\text{Circuit-SAT} \leq_P \text{SAT}$
 - $\text{IS(D)} \leq_P \text{3SAT}$

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Diagram of a reduction $X \leq_P Y$

X, Y are computational problems; R is a polynomial time transformation from input x to y so that x, y are equivalent



We used reductions

- ▶ as a means to design efficient algorithms
- ▶ for arguing about the relative hardness of problems

Decision versions of optimization problems

An optimization problem X may be transformed into a roughly equivalent problem with a **yes/no** answer, called the **decision version** $X(D)$ of the optimization problem, by

1. supplying a **target** value for the quantity to be optimized;
2. asking whether this value can be attained.

Examples:

- ▶ IS(D): given a graph G and an integer k , does G have an independent set of size k ?
- ▶ VC(D): given a graph G and an integer k , does G have a vertex cover of size k ?

The complexity class \mathcal{P} and beyond

Definition 1.

\mathcal{P} is the set of problems that can be **solved** by polynomial-time algorithms.

Beyond \mathcal{P} ?

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Beyond \mathcal{P} ?

Problems like IS(D) and VC(D):

- ▶ No polynomial time algorithm has been found despite significant effort, so we don't believe they are in \mathcal{P} .
- ▶ *Is there anything positive we can say about such problems?*

Definition 2.

An efficient certifier (or *verification algorithm*) B for a problem $X(D)$ is a **polynomial-time** algorithm that

1. takes **two** input arguments, the instance x and the *short* certificate t (both encoded as binary strings)
2. there is a polynomial $p(\cdot)$ so that for every string x , we have $x \in X(D)$ if and only if there is a string t such that $|t| \leq p(|x|)$ and $B(x, t) = \mathbf{yes}$.

Note that **existence** of the certifier B **does not** provide us with any efficient way to **solve** $X(D)$! (*why?*)

Definition 3.

We define \mathcal{NP} to be the set of decision problems that have an efficient certifier.

Fact 4.

$$\mathcal{P} \subseteq \mathcal{NP}$$

Proof.

Let $\mathbf{X(D)}$ be a problem in \mathcal{P} .

- ▶ There is an efficient algorithm A that solves $\mathbf{X(D)}$, that is, $A(x) = \mathbf{yes}$ if and only if $x \in \mathbf{X(D)}$.
- ▶ To show that $\mathbf{X(D)} \in \mathcal{NP}$, we need exhibit an efficient certifier B that takes two inputs x and t and answers \mathbf{yes} if and only if $x \in \mathbf{X(D)}$.
- ▶ The algorithm B that on inputs x, t , simply discards t and simulates $A(x)$ is such an efficient certifier.



$$\mathcal{P} = \mathcal{NP} ?$$

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- ▶ *Arguably the biggest question in theoretical CS*
- ▶ *We do not think so: finding a solution should be harder than checking one, especially for hard problems...*

Why would \mathcal{NP} contain more problems than \mathcal{P} ?

- ▶ Intuitively, the **hardest** problems in \mathcal{NP} are the **least likely** to belong to \mathcal{P} .
- ▶ How do we identify the hardest problems?

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The notion of reduction is useful again.

Definition 5 (\mathcal{NP} -complete problems:).

A problem $X(D)$ is \mathcal{NP} -complete if

1. $X(D) \in \mathcal{NP}$, and
2. for all $Y \in \mathcal{NP}$, $Y \leq_P X$.

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Fact 6.

Suppose X is \mathcal{NP} -complete. Then X is solvable in polynomial time (i.e., $X \in \mathcal{P}$) if and only if $\mathcal{P} = \mathcal{NP}$.

Why we should care whether a problem is \mathcal{NP} -complete

- ▶ If a problem is \mathcal{NP} -complete it is among the least likely to be in \mathcal{P} : it is in \mathcal{P} if and only if $\mathcal{P} = \mathcal{NP}$.

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- ▶ Therefore, from an algorithmic perspective, we need to **stop looking for efficient algorithms for the problem.**

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- ▶ Therefore, from an algorithmic perspective, we need to **stop looking for efficient algorithms for the problem.**

Instead we have a number of options

1. **approximation algorithms**, that is, algorithms that return a solution within a provable guarantee from the optimal
2. exponential algorithms practical for **small instances**
3. work on interesting **special cases**
4. study the average performance of the algorithm
5. examine **heuristics** (algorithms that work well in practice, yet provide no theoretical guarantees regarding how close the solution they find is to the optimal one)

How do we show that a problem is \mathcal{NP} -complete?

Suppose we had an \mathcal{NP} -complete problem X .

To show that another problem Y is \mathcal{NP} -complete, we only need show that

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Why?

Fact 7 (Transitivity of reductions).

If $X \leq_P Y$ and $Y \leq_P Z$, then $X \leq_P Z$.

We know that for all $A \in \mathcal{NP}$, $A \leq_P X$. By Fact 15, $A \leq_P Y$. Hence Y is \mathcal{NP} -complete.

How do we show that a problem is \mathcal{NP} -complete?

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So, *if* we had a first \mathcal{NP} -complete problem X , discovering a new problem Y in this class would require an *easier* kind of reduction: just reduce X to Y (instead of reducing *every* problem in \mathcal{NP} to Y !).

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The first \mathcal{NP} -complete problem

Theorem 7 (Cook-Levin).

Circuit SAT is \mathcal{NP} -complete.

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Syntax of Boolean expressions

- ▶ **Boolean variable** x : a variable that takes values from $\{0, 1\}$ (equivalently, $\{F, T\}$, standing for **False**, **True**).
- ▶ Suppose you are given a set of n boolean variables $\{x_1, x_2, \dots, x_n\}$.
- ▶ **Boolean connectives**: logical AND \wedge , logical OR \vee and logical NOT \neg
- ▶ **Boolean expression or Boolean formula**: boolean variables connected by boolean connectives
- ▶ **Notational convention**: ϕ is a boolean formula

Boolean expressions

A **boolean expression** may be any of the following

1. A boolean variable, e.g., x_i .
2. The **negation** of a Boolean expression ϕ , denoted by $\neg\phi_1$ or $\overline{\phi_1}$.
3. The **disjunction** (logical OR) of two Boolean expressions in parentheses ($\phi_1 \vee \phi_2$).
4. The **conjunction** (logical AND) of two Boolean expressions in parentheses ($\phi_1 \wedge \phi_2$).

Properties of Boolean expressions

Basic properties of Boolean expressions (associativity, commutativity, distribution laws)

1. $\neg\neg\phi \equiv \phi$
2. $(\phi_1 \vee \phi_2) \equiv (\phi_2 \vee \phi_1)$
3. $(\phi_1 \wedge \phi_2) \equiv (\phi_2 \wedge \phi_1)$
4. $((\phi_1 \vee \phi_2) \vee \phi_3) \equiv (\phi_1 \vee (\phi_2 \vee \phi_3))$
5. $((\phi_1 \wedge \phi_2) \wedge \phi_3) \equiv (\phi_1 \wedge (\phi_2 \wedge \phi_3))$
6. $((\phi_1 \vee \phi_2) \wedge \phi_3) \equiv ((\phi_1 \wedge \phi_3) \vee (\phi_2 \wedge \phi_3))$
7. $((\phi_1 \wedge \phi_2) \vee \phi_3) \equiv ((\phi_1 \vee \phi_3) \wedge (\phi_2 \vee \phi_3))$
8. $\neg(\phi_1 \vee \phi_2) \equiv (\neg\phi_1 \wedge \neg\phi_2)$
9. $\neg(\phi_1 \wedge \phi_2) \equiv (\neg\phi_1 \vee \neg\phi_2)$
10. $\phi_1 \vee \phi_1 \equiv \phi_1$
11. $\phi_1 \wedge \phi_1 \equiv \phi_1$

Conjunctive Normal Form (CNF)

A **literal** l_i is a variable or its negation.

Definition 8.

A Boolean formula ϕ is in CNF if it consists of **conjunctions of clauses** each of which is a **disjunction of literals**.

- ▶ In symbols, a formula ϕ with m clauses is in CNF if

$$\phi = C_1 \wedge C_2 \wedge \dots \wedge C_m$$

and each clause C_i is the disjunction of a number of literals

$$l_1 \vee l_2 \vee \dots \vee l_k$$

- ▶ **Example:** $n = 3$, $m = 2$, $\phi = (x_1 \vee \overline{x_2}) \wedge (\overline{x_1} \vee x_2 \vee x_3)$

Remark: we will henceforth work with formulas in CNF.

Semantics of boolean formulas

- ▶ Let $X = \{x_1, \dots, x_n\}$.
- ▶ A **truth assignment** for X is an assignment of truth values from $\{0, 1\}$ to each x_i .
 - ▶ So a truth assignment is a function $v : X \rightarrow \{0, 1\}$.
 - ▶ It is implied that \bar{x}_i obtains value opposite from x_i .
 - ▶ **Example:** $X = \{x_1, x_2, x_3\}$
 - ▶ Truth assignment for X : $x_1 = 1, x_2 = x_3 = 0$
- ▶ A truth assignment causes a boolean formula to receive a value from $\{0, 1\}$.
 - ▶ **Example:** $\phi = (x_1 \vee \bar{x}_2) \wedge (\bar{x}_1 \vee x_2 \vee x_3)$
 - ▶ The above truth assignment causes ϕ to evaluate to 0.

Satisfying truth assignments

- ▶ A truth assignment **satisfies a clause** if it causes the clause to evaluate to 1.
 - ▶ **Example:** $\phi = (x_1 \vee \overline{x_2}) \wedge (\overline{x_1} \vee x_2 \vee x_3)$
Then $x_1 = x_2 = 1, x_3 = 0$ satisfies both clauses in ϕ .
- ▶ A truth assignment **satisfies a formula** in CNF if it satisfies **every** clause in the formula.
 - ▶ **Example:** $x_1 = x_2 = 1, x_3 = 0$ satisfies the above ϕ .
But $x_1 = 1, x_2 = x_3 = 0$ does **not** satisfy ϕ .
- ▶ A formula ϕ is **satisfiable** if it has a satisfying truth assignment.
 - ▶ **Example:** the above ϕ is satisfiable; a *certificate* of its satisfiability is the truth assignment $x_1 = x_2 = 1, x_3 = 0$.

Satisfiability (SAT) and 3SAT

Definition 9 (SAT).

Given a formula ϕ in CNF with n variables and m clauses, is ϕ satisfiable?

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A convenient (and not easier) variant of SAT requires that **every clause** consists of **exactly three** literals.

Definition 10 (3SAT) .

Given a formula ϕ in CNF with n variables and m clauses such that each clause has exactly 3 literals, is ϕ satisfiable?

Are these problems hard?

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Theorem 11.

SAT, 3SAT are \mathcal{NP} -complete.

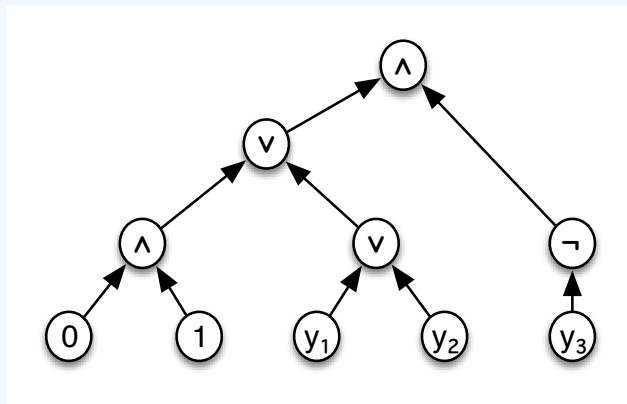
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Physical circuits and boolean combinatorial circuits

- ▶ A physical circuit consists of gates that perform logical AND, OR and NOT.
- ▶ We will model such a circuit by a **boolean combinatorial circuit** which is a labelled DAG with
 - ▶ **Source nodes**: these are the inputs of the circuit and may be hardwired to 0 or 1, or labelled with some variable.
 - ▶ **Intermediate nodes**: these correspond to the gates of the circuit and are labelled with \wedge (AND), \vee (OR) or \neg (NOT).
 - ▶ \wedge, \vee gates have two incoming and one outgoing edge
 - ▶ \neg gates have one incoming and one outgoing edge
 - ▶ **Sink node**: corresponds to the output of the circuit and has no outgoing edges.

Example circuit



A circuit C with 2 hardwired source nodes, 3 variable inputs y_1, y_2, y_3 and 5 logical gates.

Circuit-SAT: a first \mathcal{NP} -complete problem

Evaluating a circuit:

- ▶ edges are wires that carry the value of their tail node;
- ▶ intermediate nodes perform their label operation on their incoming edges, pass the result along their outgoing edge;
- ▶ the value of the circuit is the value of its output node.

Definition 12 (Circuit-SAT).

Given a circuit C , is there an assignment of truth values to its inputs that causes the output to evaluate to 1?

It is easy to see that **Circuit-SAT** is in \mathcal{NP} . Cook and Levin showed that it is \mathcal{NP} -complete.

Lemma 13.

Circuit-SAT \leq_P SAT

Intuitively, this reduction should not be too difficult: formulas and circuits are just different ways of representing boolean functions and translating from one to the other should be easy.

The following two boolean connectives are very useful.

1. $(\phi_1 \Rightarrow \phi_2)$ is a shorthand for $(\overline{\phi_1} \vee \phi_2)$.
Intuition: if $\phi_1 = 1$, then $\phi_2 = 1$ too (o.w., $(\phi_1 \Rightarrow \phi_2) = 0$).
2. $(\phi_1 \Leftrightarrow \phi_2)$ is a shorthand for $((\phi_1 \Rightarrow \phi_2) \wedge (\phi_2 \Rightarrow \phi_1))$,
which may be expanded to $(\overline{\phi_1} \vee \phi_2) \wedge (\phi_1 \vee \overline{\phi_2})$.
This clause evaluates to 1 if and only if $\phi_1 = \phi_2$.

Transforming a circuit C into a formula ϕ

Consider an arbitrary instance of **Circuit-SAT**, that is, a circuit C with source nodes, intermediate nodes and an output node.

For **every** node v in C , we introduce to ϕ

- ▶ a variable x_v that encodes the truth value computed by node v in C ;
- ▶ clauses that ensure that x_v takes on the same value as the output of node v given its inputs

Then any satisfying truth assignment for the circuit C will imply that ϕ is satisfiable, while, if ϕ is satisfiable, setting the variable inputs of C to the truth values of their corresponding variables in ϕ will result in C computing an output with value 1.

ϕ is the conjunction of the following clauses

1. If v is a source node corresponding to a variable input of the circuit C , we do not add any clause.
2. If v is a source node hardwired to 0, add $(\overline{x_v})$.
3. If v is a source node hardwired to 1, add (x_v) .
4. If v is the output node, add (x_v) .
5. If v is a node labelled by NOT and its input edge is from node u , add $(x_v \Leftrightarrow \overline{x_u})$.
6. If v is a node labelled by OR and its input edges are from nodes u and w , add $(x_v \Leftrightarrow (x_u \vee x_w))$.
7. If v is a node labelled by AND and its input edges are from nodes u and w , add $(x_v \Leftrightarrow (x_u \wedge x_w))$.

Transforming C into ϕ requires polynomial time

This completes our construction of the clauses of ϕ .

For example, for the circuit in slide 34, we construct the following formula.

$$\begin{aligned}\phi = & (\neg x_1) \wedge (x_2) \wedge (x_6 \Leftrightarrow (x_1 \wedge x_2)) \wedge (x_7 \Leftrightarrow (x_3 \vee x_4)) \wedge \\ & (x_8 \Leftrightarrow \neg x_5) \wedge (x_9 \Leftrightarrow (x_6 \vee x_7)) \wedge (x_{10} \Leftrightarrow (x_9 \wedge x_8)) \wedge (x_{10})\end{aligned}$$

The construction is polynomial in the size of the input circuit (*why?*).

Moreover, every clause consists of at most three literals, once ϕ is in CNF (*exercise*).

Proof of equivalence

- \Rightarrow Let T_C be a truth assignment to the variable inputs of C that causes C to evaluate to 1. Propagate T_C to assign a truth value to every node v in C . Define a truth assignment T_ϕ for ϕ as follows: x_v takes on the truth value of v , for every node v in C . Then T_ϕ satisfies ϕ .
- \Leftarrow Suppose ϕ has a satisfying truth assignment. Then the truth values of the variables of ϕ that correspond to inputs in C satisfy C : the clauses in ϕ guarantee that, for every node in C , the value assigned to that node is exactly what that node computes in C . Since $\phi = 1$, C evaluates to 1.

Independent set

So far, we have stated (with or without proofs) that

▶ **Circuit-SAT** is \mathcal{NP} -complete

▶ **Circuit-SAT** \leq_P **SAT**

▶ **SAT** \leq_P **3SAT**

\Rightarrow **SAT** and **3SAT** are \mathcal{NP} -complete.

Is IS(D) as “hard” as SAT?

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 - ▶ **Circuit-SAT** \leq_P **SAT**
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- \Rightarrow **SAT** and **3SAT** are \mathcal{NP} -complete.

Claim 1.

IS(D) is \mathcal{NP} -complete.

Proof.

Reduction from **3SAT**. □

Structure of the proof

Given an **arbitrary** instance formula ϕ of 3SAT, we need to transform it into a graph G and an integer k , so that

1. The transformation is completed in polynomial time.
2. The instance (G, k) is a **yes** instance of IS(D)
if and only if
 ϕ is a **yes** instance of 3SAT.

Structure of the proof

Given an **arbitrary** instance formula ϕ of 3SAT, we need to transform it into a graph G and an integer k , so that

1. The transformation is completed in polynomial time.
2. G has an independent set of size at least k
if and only if
 ϕ is satisfiable

Example: given

$$\phi = (x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee x_2 \vee \neg x_3) \wedge (x_1 \vee \neg x_2 \vee \neg x_3)$$

construct

$$(G, k)$$

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Remark 1.

- ▶ *Heart of reduction $X \leq_P Y$: understand why some **small instance** of Y makes it difficult.*
- ▶ *For IS(D), such an instance is a **triangle**: it's not clear which of its vertices to add to our independent set.*

Gadgets!

When reducing from 3SAT, we often use **gadgets**. Gadgets are constructions that ensure:

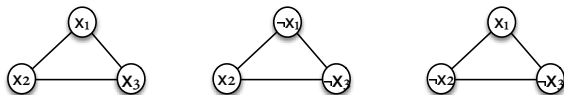
1. **Consistency of truth values in a truth assignment:** once x_i is assigned a truth value, we must henceforth consistently use it under this truth value.
2. **Clause constraints:** since ϕ is in CNF, we must provide a way to satisfy **every** clause. Equivalently, we must exhibit at least one literal that is set to 1 in every clause.

In effect, these gadgets will allow us to derive a **valid** and **satisfying** truth assignment for ϕ when the transformed instance is a **yes** instance of our problem, so we can prove equivalence of the two instances.

Gadgets for IS(D)

Clause constraint gadget: for every clause, introduce a triangle where a node is labelled by a literal in the clause.

Example: $\phi = (x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee x_2 \vee \neg x_3) \wedge (x_1 \vee \neg x_2 \vee \neg x_3)$



- ▶ Hence our graph G consists of m isolated triangles.
- ▶ The max independent set in this graph has size m : pick one vertex from every triangle. So we will set $k = m$.

Goal: derive a **truth assignment** from our independent set S .

Idea: when a node from a triangle is added to S , set the corresponding literal to 1.

2. Is this truth assignment **consistent**?
- ▶ Suppose x_1 was picked from the first triangle.
 - ▶ Can still pick \bar{x}_1 from the second triangle!
 - ▶ But then we are setting x_1 to both 1 and 0.
- ⇒ This is obviously **not** a valid truth assignment!

Consistency of truth assignment: must ensure that we cannot add a node labelled x_i **and** a node labelled \bar{x}_i to our independent set.

Consistency gadgets

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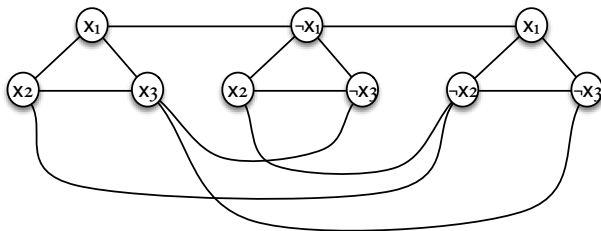
Consistency gadget: add edges between all occurrences of x_i and \bar{x}_i , for every i , in G .

Constructed instance (G, k) of IS(D)

Example: given the formula ϕ below ($n=m=3$)

$$\phi = (x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee x_2 \vee \neg x_3) \wedge (x_1 \vee \neg x_2 \vee \neg x_3),$$

the derived graph G is as follows:



Set $k=m=3$; the input instance $R(\phi)$ to IS(D) is $(G, 3)$.

Remark: the construction requires time polynomial in the size of ϕ .

Proof of equivalence

We need to show that

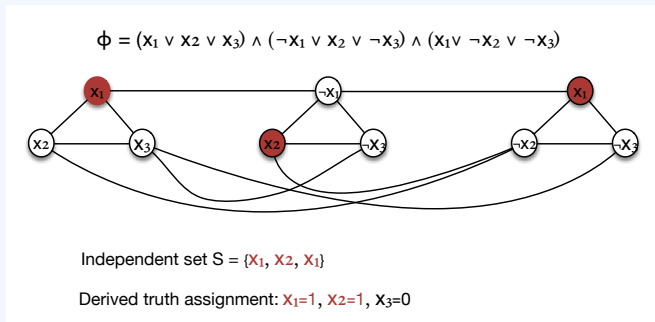
ϕ is satisfiable

if and only if

G has an independent set of size at least m

Proof of equivalence, reverse direction

- ▶ Suppose that G has an independent set S of size m .
- ▶ Then **every** triangle contributes one node to S .
- ▶ Define the following **truth assignment**
 - ▶ Set the literal corresponding to that node to 1.
 - ▶ Any variables left unset by this assignment may be set to 0 or 1 arbitrarily.



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1. **is valid**
2. **satisfies** ϕ

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We need to show that this truth assignment

1. **is valid**: by construction, x_i, \bar{x}_i cannot **both** appear in S .
2. **satisfies** ϕ : since **every** triangle contributes one node to S , every clause has a true literal, thus **every clause is satisfied**.

Proof of equivalence, forward direction

- ▶ Now suppose there is a **satisfying truth assignment** for ϕ .
- ▶ Then there is (at least) one true literal in every clause.
- ▶ Construct an **independent set** S as follows:
From every triangle, add to S a node labelled by such a literal; hence S has size m .

We claim that S thus constructed is indeed **an independent set**.

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From every triangle, add to S a node labelled by such a literal; hence S has size m .

We claim that S thus constructed is indeed **an independent set**.

1. S would not be an independent set *if* there was an edge between any two nodes in it.
2. Since all nodes in S belong to *different* triangles, an edge implies that the two nodes are labelled by opposite literals.
3. Impossible: *all* literals in S evaluate to 1.